17. A Computing Machine
17. A Computing Machine

- Overview
- Data types
- Instructions
- Operating the machine
- Machine language programming
A TOY computing machine

TOY is an imaginary machine similar to:
• Ancient computers.
• Today’s smartphone processors.
• Countless other devices designed and built over the past 50 years.

Smartphone processor, 2010s

PDP-8, 1970s
Reasons to study TOY

Prepare to learn about computer architecture
- How does your computer's processor work?
- What are its basic components?
- How do they interact?

Learn about machine-language programming.
- How do Java programs relate to your computer?
- Key to understanding Java references.
- Intellectual challenge of a new programming regime.
- Still necessary in some modern applications.

Learn fundamental abstractions that have informed processor design for decades.

multimedia, computer games, embedded devices, scientific computing,...
**Bits and words**

*Everything* in TOY is encoded with a sequence of *bits* (value 0 or 1).

- Why? Easy to represent two states (on and off) in real world.
- Bits are organized in 16-bit sequences called *words*.

More convenient for humans: *hexadecimal notation* (base 16)

- 4 *hex digits* in each word.
- Convert to and from binary 4 bits at a time.

<table>
<thead>
<tr>
<th>binary</th>
<th>hex</th>
</tr>
</thead>
<tbody>
<tr>
<td>0000</td>
<td>0</td>
</tr>
<tr>
<td>0001</td>
<td>1</td>
</tr>
<tr>
<td>0010</td>
<td>2</td>
</tr>
<tr>
<td>0011</td>
<td>3</td>
</tr>
<tr>
<td>0100</td>
<td>4</td>
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<tr>
<td>0101</td>
<td>5</td>
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<tr>
<td>0110</td>
<td>6</td>
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<tr>
<td>0111</td>
<td>7</td>
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<td>1000</td>
<td>8</td>
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<tr>
<td>1001</td>
<td>9</td>
</tr>
<tr>
<td>1010</td>
<td>A</td>
</tr>
<tr>
<td>1011</td>
<td>B</td>
</tr>
<tr>
<td>1100</td>
<td>C</td>
</tr>
<tr>
<td>1101</td>
<td>D</td>
</tr>
<tr>
<td>1110</td>
<td>E</td>
</tr>
<tr>
<td>1111</td>
<td>F</td>
</tr>
</tbody>
</table>
Inside the box

Components of TOY machine

• Memory
• Registers
• Arithmetic and logic unit (ALU)
• Program counter (PC)
• Instruction register (IR)
Memory

Holds data and instructions
- 256 words
- 16 bits in each word
- Connected to registers
- Words are *addressable*

Use *hexadecimal* for addresses
- Number words from 00 to FF
- *Think in hexadecimal*
- Use array notation
- Example: $M[2A] = \text{C024}$

Use *hexadecimal* for addresses
- Number words from 00 to FF
- *Think in hexadecimal*
- Use array notation
- Example: $M[2A] = \text{C024}$

Table of 256 words *completely specifies* contents of memory.

<table>
<thead>
<tr>
<th>Memory</th>
</tr>
</thead>
<tbody>
<tr>
<td>00</td>
</tr>
<tr>
<td>01</td>
</tr>
<tr>
<td>02</td>
</tr>
<tr>
<td>03</td>
</tr>
<tr>
<td>04</td>
</tr>
<tr>
<td>05</td>
</tr>
<tr>
<td>06</td>
</tr>
<tr>
<td>07</td>
</tr>
<tr>
<td>08</td>
</tr>
<tr>
<td>09</td>
</tr>
<tr>
<td>0A</td>
</tr>
<tr>
<td>0B</td>
</tr>
<tr>
<td>0C</td>
</tr>
<tr>
<td>0D</td>
</tr>
<tr>
<td>0E</td>
</tr>
<tr>
<td>0F</td>
</tr>
</tbody>
</table>
Arithmetic and logic unit (ALU)

ALU

- TOY’s computational engine
- A *calculator*, not a computer
- **Hardware** that implements *all* data-type operations
- How? Stay tuned for computer architecture lectures
Registers

- 16 words, addressable in hex from 0 to F (use names R[0] through R[F])
- Scratch space for calculations and data movement.
- Connected to memory and ALU
- *By convention, R[0] is always 0.*

Q. Why not just connect memory directly to ALU?
A. Too many different memory names (addresses).

Q. Why not just connect memory locations to one another?
A. Too many different connections.

Table of 16 words *completely specifies* contents of registers.
Program counter and instruction register

TOY operates by executing a sequence of instructions.

Critical abstractions in making this happen

- **Program Counter (PC)**. Memory address of next instruction.
- **Instruction Register (IR)**. Instruction being executed.

Fetch-increment-execute cycle

- Fetch: Get instruction from memory into IR.
- Increment: Update PC to point to next instruction.
- Execute: Move data to or from memory, change PC, or perform calculations, as specified by IR.
The state of the machine

Contents of memory, registers, and PC at a particular time

- Provide a record of what a program has done.
- Completely determines what the machine will do.
Image sources

http://commons.wikimedia.org/wiki/File:Calculator_casio.jpg
http://commons.wikimedia.org/wiki/File:Abacus_5.jpg
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TOY data type

A **data type** is a set of values and a set of operations on those values.

TOY’s **data type** is 16-bit *two's complement* integers.

Two kinds of operations

- Arithmetic.
- Bitwise.

All other types of data must be implemented with **software**

- 32-bit and 64-bit integers.
- 32-bit and 64-bit floating point values.
- Characters and strings.
- ...

*All values are represented in 16-bit words.*
TOY data type (original design): Unsigned integers

Values. 0 to $2^{16} - 1$, encoded in binary (or, equivalently, hex).

Example. $6375_{10}$.

<table>
<thead>
<tr>
<th>binary</th>
<th>15</th>
<th>14</th>
<th>13</th>
<th>12</th>
<th>11</th>
<th>10</th>
<th>9</th>
<th>8</th>
<th>7</th>
<th>6</th>
<th>5</th>
<th>4</th>
<th>3</th>
<th>2</th>
<th>1</th>
<th>0</th>
</tr>
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<tbody>
<tr>
<td></td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>1</td>
<td>1</td>
<td>0</td>
<td>0</td>
<td>1</td>
<td>1</td>
<td>1</td>
<td>0</td>
<td>0</td>
<td>1</td>
<td>1</td>
<td>1</td>
<td></td>
</tr>
</tbody>
</table>

$2^{12} + 2^{11}$

$+2^7 + 2^6 + 2^5$

$+2^2 + 2^1 + 2^0$

hex

<table>
<thead>
<tr>
<th>1 $\times 16^3$</th>
<th>8 $\times 16^2$</th>
<th>E $\times 16^1$</th>
<th>7</th>
</tr>
</thead>
<tbody>
<tr>
<td>4096</td>
<td>2048</td>
<td>224</td>
<td>7</td>
</tr>
</tbody>
</table>

Operations.
- Add.
- Subtract.
- Test if 0.

Example. $18E7 + 18E7 = 31CE$

\[
\begin{array}{cccccccccccccccc}
0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 1 & 1 \\
+ & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 1 & 1 \\
= & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 1 & 1 & 0 \\
\end{array}
\]

Warning. TOY ignores overflow.
TOY data type (better design): two's complement

Values.  – $2^{15}$ to $2^{15} - 1$, encoded in 16-bit two's complement.

Operations.
• Add.
• Subtract.
• Test if positive, negative, or 0.

16 bit two's complement
• 16-bit binary representation of x for positive x.
• 16-bit binary representation of $2^{16} - |x|$ for negative x.

Useful properties
• Leading bit (bit 15) signifies sign.
• 0000000000000000 represents zero.
• Add/subtract is the same as for unsigned.

<table>
<thead>
<tr>
<th>decimal</th>
<th>hex</th>
<th>binary</th>
</tr>
</thead>
<tbody>
<tr>
<td>+32,767</td>
<td>7FF</td>
<td>0111111111111111</td>
</tr>
<tr>
<td>+32,766</td>
<td>7FE</td>
<td>0111111111111110</td>
</tr>
<tr>
<td>+32,765</td>
<td>7FD</td>
<td>0111111111111111</td>
</tr>
<tr>
<td>...</td>
<td></td>
<td></td>
</tr>
<tr>
<td>+3</td>
<td>0003</td>
<td>0000000000000011</td>
</tr>
<tr>
<td>+2</td>
<td>0002</td>
<td>0000000000000010</td>
</tr>
<tr>
<td>+1</td>
<td>0001</td>
<td>0000000000000001</td>
</tr>
<tr>
<td>0</td>
<td>0000</td>
<td>0000000000000000</td>
</tr>
<tr>
<td>−1</td>
<td>FFFF</td>
<td>1111111111111111</td>
</tr>
<tr>
<td>−2</td>
<td>FFEE</td>
<td>1111111111111110</td>
</tr>
<tr>
<td>−3</td>
<td>FFDD</td>
<td>1111111111111101</td>
</tr>
<tr>
<td>...</td>
<td></td>
<td></td>
</tr>
<tr>
<td>−32,766</td>
<td>8002</td>
<td>1000000000000000</td>
</tr>
<tr>
<td>−32,767</td>
<td>8001</td>
<td>1000000000000000</td>
</tr>
<tr>
<td>−32,768</td>
<td>8000</td>
<td>1000000000000000</td>
</tr>
</tbody>
</table>

slight annoyance: one extra negative value
Two's complement: conversion

To convert from decimal to two's complement
- If greater than +32,767 or less than −32,768 report error.
- Convert to 16-bit binary.
- If not negative, done.
- If negative, flip all bits and add 1.

To convert from two's complement to decimal
- If sign bit is 1, flip all bits and add 1 and output minus sign.
- Convert to decimal.

To add/subtract
- Use same rules as for unsigned binary.
- (Still) ignore overflow.

Examples

<table>
<thead>
<tr>
<th>decimal</th>
<th>binary</th>
<th>hex</th>
</tr>
</thead>
<tbody>
<tr>
<td>+13</td>
<td>0000000000001101</td>
<td>000D</td>
</tr>
<tr>
<td>−13</td>
<td>1111111111110011</td>
<td>FFF3</td>
</tr>
<tr>
<td>+256</td>
<td>0000000100000000</td>
<td>0100</td>
</tr>
<tr>
<td>−256</td>
<td>1111111100000000</td>
<td>FF00</td>
</tr>
</tbody>
</table>

Examples

<table>
<thead>
<tr>
<th>binary</th>
<th>decimal</th>
<th>hex</th>
</tr>
</thead>
<tbody>
<tr>
<td>0001</td>
<td>0000000000000001</td>
<td>110</td>
</tr>
<tr>
<td>FFFF</td>
<td>1111111111111111</td>
<td>−110</td>
</tr>
<tr>
<td>FF0D</td>
<td>1111111100001101</td>
<td>−24310</td>
</tr>
<tr>
<td>00F3</td>
<td>0000000011110011</td>
<td>+24310</td>
</tr>
</tbody>
</table>

Example

<table>
<thead>
<tr>
<th>binary</th>
<th>decimal</th>
<th>hex</th>
</tr>
</thead>
<tbody>
<tr>
<td>−256</td>
<td>1111111100000000</td>
<td>FF00</td>
</tr>
<tr>
<td>+13</td>
<td>+0000000000001011</td>
<td>+000D</td>
</tr>
<tr>
<td>= −243</td>
<td>=1111111100001101</td>
<td>=FF0D</td>
</tr>
</tbody>
</table>
Overflow in two's complement

\[
\begin{array}{ccc}
32,767_{10} = 2^{15} - 1 & 0111111111111111 & 7FFF \\
+1 & +0000000000000001 & +0001 \\
= 1000000000000000 & = 8000 & = -2^{15} = -32,768_{10}
\end{array}
\]

largest (positive) number

smallest (negative) number

http://xkcd.com/571/
TOY data type: Bitwise operations

Operations
- Bitwise AND.
- Bitwise XOR.
- Shift left.
- Shift right.

Special note: Shift left/right operations also implement multiply/divide by powers of 2 for integers.
Image sources

https://xkcd.com/571/
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### TOY instructions

**ANY** 16-bit (4 hex digit) value defines a TOY instruction.

**First hex digit** specifies which instruction.

**Each instruction changes machine state in a well-defined way.**

<table>
<thead>
<tr>
<th>category</th>
<th>opcodes</th>
<th>implements</th>
<th>changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>operations</td>
<td>1 2 3 4 5 6</td>
<td>data-type operations</td>
<td>registers</td>
</tr>
<tr>
<td>data movement</td>
<td>7 8 9 A B</td>
<td>data moves between registers and memory</td>
<td>registers, memory</td>
</tr>
<tr>
<td>flow of control</td>
<td>0 C D E F</td>
<td>conditionals, loops, and functions</td>
<td>PC</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>opcode</th>
<th>instruction</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>halt</td>
</tr>
<tr>
<td>1</td>
<td>add</td>
</tr>
<tr>
<td>2</td>
<td>subtract</td>
</tr>
<tr>
<td>3</td>
<td>bitwise and</td>
</tr>
<tr>
<td>4</td>
<td>bitwise xor</td>
</tr>
<tr>
<td>5</td>
<td>shift left</td>
</tr>
<tr>
<td>6</td>
<td>shift right</td>
</tr>
<tr>
<td>7</td>
<td>load address</td>
</tr>
<tr>
<td>8</td>
<td>load</td>
</tr>
<tr>
<td>9</td>
<td>store</td>
</tr>
<tr>
<td>A</td>
<td>load indirect</td>
</tr>
<tr>
<td>B</td>
<td>store indirect</td>
</tr>
<tr>
<td>C</td>
<td>branch if zero</td>
</tr>
<tr>
<td>D</td>
<td>branch if positive</td>
</tr>
<tr>
<td>E</td>
<td>jump register</td>
</tr>
<tr>
<td>F</td>
<td>jump and link</td>
</tr>
</tbody>
</table>
Encoding instructions

ANY 16-bit (4 hex digit) value defines a TOY instruction.

Two different instruction formats

• **Type RR**: Opcode and 3 registers.
  
<table>
<thead>
<tr>
<th>15</th>
<th>14</th>
<th>13</th>
<th>12</th>
<th>11</th>
<th>10</th>
<th>9</th>
<th>8</th>
<th>7</th>
<th>6</th>
<th>5</th>
<th>4</th>
<th>3</th>
<th>2</th>
<th>1</th>
<th>0</th>
</tr>
</thead>
<tbody>
<tr>
<td>opcode</td>
<td>destination d</td>
<td>source s</td>
<td>source t</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
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</tr>
</tbody>
</table>

• **Type A**: Opcode, 1 register, and 1 memory address.
  
<table>
<thead>
<tr>
<th>15</th>
<th>14</th>
<th>13</th>
<th>12</th>
<th>11</th>
<th>10</th>
<th>9</th>
<th>8</th>
<th>7</th>
<th>6</th>
<th>5</th>
<th>4</th>
<th>3</th>
<th>2</th>
<th>1</th>
<th>0</th>
</tr>
</thead>
<tbody>
<tr>
<td>opcode</td>
<td>destination d</td>
<td>address addr</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
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</tbody>
</table>

Examples

<table>
<thead>
<tr>
<th>opcode</th>
<th>instruction</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>RR halt</td>
</tr>
<tr>
<td>1</td>
<td>RR add</td>
</tr>
<tr>
<td>2</td>
<td>RR subtract</td>
</tr>
<tr>
<td>3</td>
<td>RR bitwise and</td>
</tr>
<tr>
<td>4</td>
<td>RR bitwise xor</td>
</tr>
<tr>
<td>5</td>
<td>RR shift left</td>
</tr>
<tr>
<td>6</td>
<td>RR shift right</td>
</tr>
<tr>
<td>7</td>
<td>A load address</td>
</tr>
<tr>
<td>8</td>
<td>A load</td>
</tr>
<tr>
<td>9</td>
<td>A store</td>
</tr>
<tr>
<td>A</td>
<td>RR load indirect</td>
</tr>
<tr>
<td>B</td>
<td>RR store indirect</td>
</tr>
<tr>
<td>C</td>
<td>A branch if zero</td>
</tr>
<tr>
<td>D</td>
<td>A branch if positive</td>
</tr>
<tr>
<td>E</td>
<td>RR jump register</td>
</tr>
<tr>
<td>F</td>
<td>A jump and link</td>
</tr>
</tbody>
</table>

1CAB \ Add R[A] to R[B] and put result in R[C].

8A15 \ Load into R[A] data from M[15].
Add two integers
• Load operands from memory into registers.
• Add the registers.
• Put result in memory.

**Load** into R[A] data from M[15]
**Load** into R[B] data from M[16]
**Add** R[A] and R[B] and put result into R[C]
**Store** R[C] into M[17]
**Halt**

**Q.** How can you tell whether a word is an instruction?

**A.** If the PC has its address, it *is* an instruction!
Same program with different data

Add two integers
• Load operands from memory into registers.
• Add the registers.
• Put result in memory.

Load into R[A] data from M[15]
Load into R[B] data from M[16]
Add R[A] and R[B] and put result into R[C]
Store R[C] into M[17]
Halt

Q. How can you tell whether a word is data?

A. If it is added to another word, it is data!
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Outside the box

User interface
• Switches.
• Lights.
• Control Buttons.

First step: Turn on the machine!
To load a program into memory

To load an instruction

- Set 8 memory address switches.
- Set 16 data switches to instruction encoding.
- Press LOAD to load instruction from switches into addressed memory word.

LOAD

LOOK

STEP

RUN

ON

LOAD

A COMPUTING MACHINE

ADDR

DATA

10: 8A15
11: 8B16
12: 1CAB
13: 9C17
14: 0000
To load an instruction

- Set 8 memory address switches.
- Set 16 data switches to instruction encoding.
- Press LOAD to load instruction from switches into addressed memory word.
Loading instructions into memory

To load an instruction
- Set 8 memory address switches.
- Set 16 data switches to instruction encoding.
- Press LOAD to load instruction from switches into addressed memory word.
Loading instructions into memory

To load an instruction

- Set 8 memory address switches.
- Set 16 data switches to instruction encoding.
- Press LOAD to load instruction from switches into addressed memory word.
Loading instructions into memory

**To load an instruction**
- Set 8 memory address switches.
- Set 16 data switches to instruction encoding.
- Press LOAD to load instruction from switches into addressed memory word.
Looking at what's in the memory

To double check that you loaded the data correctly
- Set 8 memory address switches.
- Press LOOK to examine the addressed memory word.

<table>
<thead>
<tr>
<th>ADDR</th>
<th>DATA</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
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</tr>
<tr>
<td></td>
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</tr>
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<td></td>
<td></td>
</tr>
</tbody>
</table>

LOAD  LOOK  STEP  RUN  ON

10: 8A15
11: 8B16
12: 1CAB
13: 9C17
14: 0000
Loading data into memory

To load data, use the same procedure as for instructions

- Set 8 memory address switches.
- Set 16 data switches to data encoding.
- Press LOAD to load data from switches into addressed memory word.

LOAD
LOOK
STEP
RUN
ON

TOY
A COMPUTING MACHINE

ADDR

DATA

10: 8A15
11: 8B16
12: 1CAB
13: 9C17
14: 0000

15: 0008
16: 0005
Loading data into memory

To load data, use the same procedure as for instructions

- Set 8 memory address switches.
- Set 16 data switches to data encoding.
- Press LOAD to load data from switches into addressed memory word.

LOAD  LOOK  STEP  RUN  ON

10: 8A15
11: 8B16
12: 1CAB
13: 9C17
14: 0000

15: 0008
16: 0005
Running a program

To run a program, set the address switches to the address of first instruction and press RUN.

[ data lights may flash, but all (and RUN light) go off when HALT instruction is reached ]

To see the output, set the address switches to the address of expected result and press LOOK.

To run the program again, enter different data and press RUN again.
Q. Did people really program this way?

A. Yes!
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Machine language programming

TOY instructions support the same **basic programming constructs** as Java.

- Primitive data types.
- Assignment statements.
- Conditionals and loops.
- Arrays (next lecture).
- Standard input and output (next).

*and* can support advanced programming constructs, as well.

- Functions and libraries (see text).
- Linked structures (see text).
### Conditionals and loops

To control the flow of instruction execution

- Test a register's value.
- Change the PC, depending on the value.

**Example: Absolute value of R[A]**

<table>
<thead>
<tr>
<th>Opcode</th>
<th>Instruction</th>
</tr>
</thead>
<tbody>
<tr>
<td>DA12</td>
<td>If R[A] &gt; 0 set PC to 12 (skip 11)</td>
</tr>
<tr>
<td>2A0A</td>
<td>Subtract R[A] from 0 (R[0]) and put result into R[A]</td>
</tr>
</tbody>
</table>

**Example: Typical while loop (assumes R[1] is 0001)**

```plaintext
while (a != 0) {
    ...
    a--;
}
```

<table>
<thead>
<tr>
<th>Opcode</th>
<th>Instruction</th>
</tr>
</thead>
<tbody>
<tr>
<td>CA15</td>
<td>If R[A] is 0 set PC to 15</td>
</tr>
<tr>
<td>2AA1</td>
<td>Decrement R[A] by 1</td>
</tr>
<tr>
<td>C010</td>
<td>Set PC to 10</td>
</tr>
</tbody>
</table>

**opcode**  **instruction**

| C          | branch if zero                  |
| D          | branch if positive              |
An immediate problem

• We're not going to be able to address real-world problems with just switches and lights for I/O!
• One solution: Paper tape.

Need to bolt new I/O devices to the side of the machine.
### Standard input and output

#### Punched paper tape
- Encode 16-bit words in two 8-bit rows.
- To *write* a word, *punch a hole* for each 1.
- To *read* a word, shine a light behind the tape and sense the holes.

#### TOY mechanism
- Connect hardware to memory location FF.
- To *write* the contents of a register to stdout, *store* to FF.
- To *read* from stdin into a register, *load* from FF.
Flow of control and standard output example: Fibonacci numbers

| Register trace | A   | 1 | 1 | 2 | 3 | 5 | 8 | 13 | 21 | 34 | 55 | 89 |
|                | B   | 1 | 2 | 3 | 5 | 8 | 13 | 21 | 34 | 55 | 89 | 144 |
|                | C   | 2 | 3 | 3 | 5 | 8 | 13 | 21 | 34 | 55 | 89 | 144 |
|                | 9   | A | 9 | 8 | 7 | 6 | 5 | 4  | 3  | 2  | 1  | 0  |

```
40 7 1 0 1  R[1] = 1
41 7 A 0 1  R[A] = 1
42 7 B 0 1  R[B] = 1
43 8 9 4 C  R[9] = M[4C]
44 C 9 4 B  if (R[9] == 0) PC = 4B
45 9 A F F  write R[A] to stdout
46 1 C A B  R[C] = R[A] + R[B]
47 1 A B 0  R[A] = R[B]
48 1 B C 0  R[B] = R[C]
4A C 0 4 4  PC = 44
4B 0 0 0 0  halt
4C 0 0 0 A
```

a = 1;
b = 1;
i = N;
while (i != 0) {
    StdOut.print(a);
    c = a + b;
a = b;
b = c;
i = i - 1;
}
### TOY reference card

<table>
<thead>
<tr>
<th>opcode</th>
<th>operation</th>
<th>format</th>
<th>pseudo-code</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>halt</td>
<td>—</td>
<td>halt</td>
</tr>
<tr>
<td>1</td>
<td>add</td>
<td>RR</td>
<td>R[d] = R[s] + R[t]</td>
</tr>
<tr>
<td>2</td>
<td>subtract</td>
<td>RR</td>
<td>R[d] = R[s] - R[t]</td>
</tr>
<tr>
<td>3</td>
<td>bitwise and</td>
<td>RR</td>
<td>R[d] = R[s] &amp; R[t]</td>
</tr>
<tr>
<td>4</td>
<td>bitwise xor</td>
<td>RR</td>
<td>R[d] = R[s] ^ R[t]</td>
</tr>
<tr>
<td>5</td>
<td>shift left</td>
<td>RR</td>
<td>R[d] = R[s] &lt;&lt; R[t]</td>
</tr>
<tr>
<td>6</td>
<td>shift right</td>
<td>RR</td>
<td>R[d] = R[s] &gt;&gt; R[t]</td>
</tr>
<tr>
<td>7</td>
<td>load addr</td>
<td>A</td>
<td>R[d] = addr</td>
</tr>
<tr>
<td>8</td>
<td>load</td>
<td>A</td>
<td>R[d] = M[addr]</td>
</tr>
<tr>
<td>9</td>
<td>store</td>
<td>A</td>
<td>M[addr] = R[d]</td>
</tr>
<tr>
<td>A</td>
<td>load indirect</td>
<td>RR</td>
<td>R[d] = M[R[t]]</td>
</tr>
<tr>
<td>B</td>
<td>store indirect</td>
<td>RR</td>
<td>M[R[t]] = R[d]</td>
</tr>
<tr>
<td>C</td>
<td>branch zero</td>
<td>A</td>
<td>if (R[d] == 0) PC = addr</td>
</tr>
<tr>
<td>D</td>
<td>branch positive</td>
<td>A</td>
<td>if (R[d] &gt; 0) PC = addr</td>
</tr>
<tr>
<td>E</td>
<td>jump register</td>
<td>RR</td>
<td>PC = R[d]</td>
</tr>
<tr>
<td>F</td>
<td>jump and link</td>
<td>A</td>
<td>R[d] = PC + 1; PC = addr</td>
</tr>
</tbody>
</table>

**Format RR**

```
15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0
```

- **opcode**
- **destination d**
- **source s**
- **source t**

**Format A**

```
15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0
```

- **opcode**
- **destination d**
- **address ADDR**

**ZERO** R[0] is always 0.

**STANDARD INPUT** Load from FF.

**STANDARD OUTPUT** Store to FF.
Pop quiz 1 on TOY

Q. What is the interpretation of

1A75 as a TOY instruction?

1A75 as a two's complement integer value?

0FFF as a TOY instruction?

0FFF as a two's complement integer value?

8888 as a TOY instruction?

8888 as a two's complement integer value? (Answer in base 16).
Q. How does one flip all the bits in a TOY register?
Q. What does the following TOY program leave in $R[2]$?

<p>| | | | | | | |</p>
<table>
<thead>
<tr>
<th></th>
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</tr>
</thead>
<tbody>
<tr>
<td>10</td>
<td>7</td>
<td>C</td>
<td>0</td>
<td>A</td>
<td></td>
<td></td>
</tr>
<tr>
<td>11</td>
<td>7</td>
<td>1</td>
<td>0</td>
<td>1</td>
<td></td>
<td></td>
</tr>
<tr>
<td>12</td>
<td>7</td>
<td>2</td>
<td>0</td>
<td>1</td>
<td></td>
<td></td>
</tr>
<tr>
<td>13</td>
<td>1</td>
<td>2</td>
<td>2</td>
<td>2</td>
<td></td>
<td></td>
</tr>
<tr>
<td>14</td>
<td>2</td>
<td>C</td>
<td>C</td>
<td>1</td>
<td></td>
<td></td>
</tr>
<tr>
<td>15</td>
<td>D</td>
<td>C</td>
<td>1</td>
<td>3</td>
<td></td>
<td></td>
</tr>
<tr>
<td>16</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

$R[C] = 10_{10}$

$R[1] = 1$

$R[2] = 1$


$R[C] = R[C] - 1$

if $(R[C] > 0)$ PC = 13

HALT
TOY vs. your laptop

Two different computing machines

• Both implement basic data types, conditionals, loops, and other low-level constructs.
• Both can have arrays, functions, and other high-level constructs.
• Both have infinite input and output streams.

Q. Is 256 words enough to do anything useful?
A. Yes! (See book, and stay tuned for next lecture.)

A. Yes! It is a Turing Machine, with a read/write I/O device (see theory lectures).

OK, we definitely want a faster version with more memory when we can afford it...
17. A Computing Machine